OPERATING SYSTEMS

LABORATORY MANUAL

DEPARTMENT OF COMPUTER SCIENCE AND ENGINEERING

DEPARTMENT OF COMPUTER SCIENCE AND ENGINEERING

VISION

To improve the quality of technical education that provides efficient software engineers with an attitude to adapt challenging IT needs of local, national and international arena, through teaching and interaction with alumni and industry.

MISSION

Department intends to meet the contemporary challenges in the field of IT and is playing a vital role in shaping the education of the 21st century by providing unique educational and research opportunities.

PROGRAMME EDUCATIONAL OBJECTIVES (PEOs)

PEO1 - ANALYTICAL SKILLS

To facilitate the graduates with the ability to visualize, gather information, articulate, analyze, solve complex problems, and make decisions. These are essential to address the challenges of complex and computation intensive problems increasing their productivity.

PEO2 - TECHNICAL SKILLS

To facilitate the graduates with the technical skills that prepare them for immediate employment and pursue certification providing a deeper understanding of the technology in advanced areas of computer science and related fields, thus encouraging to pursue higher education and research based on their interest.

PEO3 - SOFT SKILLS

To facilitate the graduates with the soft skills that include fulfilling the mission, setting goals, showing self-confidence by communicating effectively, having a positive attitude, get involved in team- work, being a leader, managing their career and their life.

PEO4 - PROFESSIONAL ETHICS

To facilitate the graduates with the knowledge of professional and ethical responsibilities by paying attention to grooming, being conservative with style, following dress codes, safety codes, and adapting themselves to technological advancements.

PROGRAM SPECIFIC OUTCOMES (PSOs)

After the completion of the course, B. Tech Information Technology, the graduates will have the following Program Specific Outcomes:

- 1. Fundamentals and critical knowledge of the Computer System: Able to Understand the working principles of the computer System and its components , Apply the knowledge to build, asses, and analyze the software and hardware aspects of it .
- 2. The comprehensive and Applicative knowledge of Software Development: Comprehensive skills of Programming Languages, Software process models, methodologies, and able to plan, develop, test, analyze, and manage the software and hardware intensive systems in heterogeneous platforms individually or working in teams.
- 3. Applications of Computing Domain & Research: Able to use the professional, managerial, interdisciplinary skill set, and domain specific tools in development processes, identify the research gaps, and provide innovative solutions to them.

PROGRAMOUTCOMES (POs)

Engineering Graduates should possess the following:

- 1. **Engineering knowledge:** Apply the knowledge of mathematics, science, engineering fundamentals, and an engineering specialization to the solution of complex engineering problems.
- 2. **Problem analysis:** Identify, formulate, review research literature, and analyze complex engineering problems reaching substantiated conclusions using first principles of mathematics, natural sciences, and engineering sciences.
- 3. **Design / development of solutions:** Design solutions for complex engineering problems and design system components or processes that meet the specified needs with appropriate consideration for the public health and safety, and the cultural, societal, and environmental considerations.
- 4. Conduct investigations of complex problems: Use research-based knowledge and research methods including design of experiments, analysis and interpretation of data, and synthesis of the information to provide valid conclusions.
- 5. Modern tool usage: Create, select, and apply appropriate techniques, resources, and modern engineering and IT tools including prediction and modeling to complex engineering activities with an understanding of the limitations.
- 6. The engineer and society: Apply reasoning informed by the contextual knowledge to assess societal, health, safety, legal and cultural issues and the consequent responsibilities relevant to the professional engineering practice.
- 7. **Environment and sustainability:** Understand the impact of the professional engineering solutions in societal and environmental contexts, and demonstrate the knowledge of, and need for sustainable development.
- 8. **Ethics:** Apply ethical principles and commit to professional ethics and responsibilities and norms of the engineering practice.
- 9. **Individual and team work:** Function effectively as an individual, and as a member or leader in diverse teams, and in multidisciplinary settings.
- 10. **Communication:** Communicate effectively on complex engineering activities with the engineering community and with society at large, such as, being able to comprehend and write effective reports and design documentation, make effective presentations, and give and receive clear instructions.
- 11. **Project management and finance**: Demonstrate knowledge and understanding of the engineering and management principles and apply these to one's own work, as a member and leader in a team, to manage projects and in multi-disciplinary environments.
- 12. **Life- long learning**: Recognize the need for, and have the preparation and ability to engage in independent and life-long learning in the broadest context of technological change.

OBJECTIVES:

- To understand the functionalities of various layers of OSI model
- To explain the difference between hardware, software; operating systems, programs and files.
 - Identify the purpose of different software applications.

Week 1: Simulate the following CPU scheduling algorithms.

a) FCFS

- b) SJF
- c) Round Robin
- d) Priority.

Week 2: Write a C program to simulate producer-consumer problem using Semaphores

Week 3: Write a C program to simulate the concept of Dining-philosophers problem.

Week 4: Simulate MVT and MFT.

Week 5: Write a C program to simulate the following contiguous memory allocation Techniques

a) Worst fit b) Best fit

- c) First fit.

Week 6: Simulate all page replacement algorithms a)FIFO b) LRU c) OPTIMAL

Week 7: Simulate all File Organization Techniques Single level directory b) Two level directory

Week 8: Simulate all file allocation strategies Sequential b) Indexed c) Linked.

Week 9: Simulate Bankers Algorithm for Dead Lock Avoidance. Week 10: Simulate Bankers Algorithm for Dead Lock Prevention. Week 11: Write a C program to simulate disk scheduling algorithms.

a) FCFS

a)

b) SCAN

c) C-SCAN

REFERENCE BOOKS:

- An Introduction to Operating Systems, P.C.P Bhatt, 2nd edition, PHI. 1.
- 2. Modern Operating Systems, Andrew S Tanenbaum, 3rd Edition, PHI **OUTCOMES:**

At the end of the course the students are able to:

- Ability to implement inter process communication between two processes.
- Ability to design and solve synchronization problems.
- Ability to simulate and implement operating system concepts such as scheduling, Deadlock management, file management, and memory management.

THE NEOTIA UNIVERSITY

DEPARTMENT OF COMPUTER SCIENCE AND ENGINEERING

Operating Systems Lab Manual

TABLE OF CONTENTS

EXP.NO	NAMEOF THE EXPERIMENT	PAGE.NO		
EZXINO	CPU SCHEDULING ALGORITHMS	TAGE		
	Careconic age was scanningstiminated orange to the west-decision as a consequence	4.7		
1	A) FIRST COME FIRST SERVE(FCFS)	1-3		
	B) SHORTEST JOB FIRST(SJF)	4-6		
	C) ROUND ROBIN	7-9		
	D) PRIORITY	10-12		
2	PRODUCER-CONSUMER PROBLEM USING	13-14		
	SEMAPHORES 13-14			
3	DINING-PHILOSOPHERS PROBLEM	15-18		
	MEMORYMANAGEMENT TECHNIQUE	ES		
	A) MULTI PROGRAMMING WITH FIXED	10.21		
4	NUMBER OF TASKS(MFT)	19-21		
	B) MULTI PROGRAMMING WITH	22.24		
	VARIABLE NUMBER OF TASKS(MVT)	22-24		
	CONTIGUOUS MEMORY ALLOCATIO	N		
_	A) WORST FIT	25-26		
5	B) BEST FIT	27-28		
	C) FIRST FIT	28-29		
PAGE REPLACEMENT ALGORITHMS				
_	A) FIRST IN FIRST OUT(FIFO)	30-32		
6	B) LEAST RECENTLY USED(LRU)	33-35		
	C) OPTIMAL	36-39		
	FILE ORGANIZATION TECHNIQUES			
7	A) SINGLE LEVEL DIRECTORY	40-42		
	B) TWO LEVEL DIRECTORY	43-46		
	FILE ALLOCATION STRATEGIES	CHR025 92-50-06		
	A) SEQUENTIAL	47-49		
8	B) INDEXED	50-52		
	C) LINKED	53-55		
9	DEAD LOCK AVOIDANCE	56-59		
10	DEAD LOCK PREVENTION	60-62		
	DISK SCHEDULING ALGORITHMS	Constitution and Security		
	A) FCFS	63-64		
11	B) SCAN	65-66		
	C) C-SCAN	67-69		

CPU SCHEDULINGALGORITHMS

A). FIRST COME FIRST SERVE:

AIM: To write a c program to simulate the CPU scheduling algorithm First Come First Serve (FCFS)

DESCRIPTION:

To calculate the average waiting time using the FCFS algorithm first the waiting time of the first process is kept zero and the waiting time of the second process is the burst time of the first process and the waiting time of the third process is the sum of the burst times of the first and the second process and so on. After calculating all the waiting times the average waiting time is calculated as the average of all the waiting times. FCFS mainly says first come first serve the algorithm which came first will be served first.

ALGORITHM:

- Step 1: Start the process
- Step 2: Accept the number of processes in the ready Queue
- Step 3: For each process in the ready Q, assign the process name and the burst time Step
- 4: Set the waiting of the first process as _0'and its burst time as its turnaround time Step
- 5: for each process in the Ready Q calculate
- a). Waiting time (n) = waiting time (n-1) + Burst time (n-1) b).

Turnaround time (n)= waiting time(n)+Burst time(n)

- Step 6: Calculate
 - a) Average waiting time = Total waiting Time / Number of process
- b) Average Turnaround time = Total Turnaround Time / Number of

process Step 7: Stop the process

SOURCE CODE:

```
#include<stdio.h>
#include<conio.h>
main()
1
int bt[20], wt[20], tat[20], i, n;
float wtavg, tatavg;
clrscr();
printf("\nEnter the number of processes
"); scanf("%d", &n);
for(i=0;i<n;i++)
{
printf("\nEnter Burst Time for Process %d -- ", i);
scanf("%d", &bt[i]);
wt[0] = wtavg = 0;
tat[0] = tatavg = bt[0];
for(i=1;i<n;i++)
{
wt[i] = wt[i-1] + bt[i-1];
tat[i] = tat[i-1] +bt[i];
wtavg = wtavg + wt[i];
tatavg = tatavg + tat[i];
printf("\t PROCESS \tBURST TIME \t WAITING TIME\t TURNAROUND TIME\n");
for(i=0;i<n;i++)
       printf("\n\t P%d \t\t %d \t\t %d \t\t %d", i, bt[i], wt[i], tat[i]);
       printf("\nAverage Waiting Time -- %f", wtavg/n);
printf("\nAverage Turnaround Time -- %f", tatavg/n);
getch();
}
```

INPUT

Enter the number of processes	3
Enter Burst Time for Process 0	24
Enter Burst Time for Process 1	3
Enter Burst Time for Process 2	3

OUTPUT

PROCESS	BURST TIME	WAITING TIME	TURNAROUN D
			TIME
P0	24	0	24
P1	3	24	27
P2	3	27	30
Average Waiting Time	e 17.000000		
Average Turnaround	Гime	27.000000	

B). SHORTEST JOB FIRST:

AIM: To write a program to stimulate the CPU scheduling algorithm Shortest job first (Non- Preemption)

DESCRIPTION:

To calculate the average waiting time in the shortest job first algorithm the sorting of the process based on their burst time in ascending order then calculate the waiting time of each process as the sum of the bursting times of all the process previous or before to that process.

ALGORITHM:

- Step 1: Start the process
- Step 2: Accept the number of processes in the ready Queue
- Step 3: For each process in the ready Q, assign the process id and accept the CPU burst time
- Step 4: Start the Ready Q according the shortest Burst time by sorting according to lowest to highest burst time.
- Step 5: Set the waiting time of the first process as _0' and its turnaround time as its burst time.
- Step 6: Sort the processes names based on their Burt time
- Step 7: For each process in the ready queue, calculate
- a) Waiting time (n-1) + Burst time (n-1)
- b) Turnaround time (n)= waiting time(n)+Burst time(n)
- Step 8: Calculate
 - c) Average waiting time = Total waiting Time / Number of process
- d) Average Turnaround time = Total Turnaround Time / Number of process Step 9: Stop the process

SOURCE CODE:

```
#include<stdio.h>
#include<conio.h>
main()
1
int p[20], bt[20], wt[20], tat[20], i, k, n, temp; float wtavg,
clrscr();
printf("\nEnter the number of processes -- ");
scanf("%d", &n);
for(i=0;i<n;i++)
{
p[i]=i;
printf("Enter Burst Time for Process %d -- ", i);
scanf("%d", &bt[i]);
for(i=0;i \le n;i++)
for(k=i+1;k\leq n;k++)
if(bt[i]>bt[k])
temp=bt[i];
bt[i]=bt[k];
bt[k]=temp;
temp=p[i];
p[i]=p[k];
p[k]=temp;
wt[0] = wtavg = 0;
tat[0] = tatavg = bt[0]; for(i=1;i \le n;i++)
wt[i] = wt[i-1] + bt[i-1];
tat[i] = tat[i-1] + bt[i];
wtavg = wtavg + wt[i];
tatavg = tatavg + tat[i];
}
printf("\n\t PROCESS \tBURST TIME \t WAITING TIME\t TURNAROUND TIME\n");
for(i=0;i< n;i++)
       printf("\n\t P%d \t\t %d \t\t %d", p[i], bt[i], wt[i], tat[i]);
       printf("\nAverage Waiting Time -- %f", wtavg/n);
printf("\nAverage Turnaround Time -- %f", tatavg/n); getch();}
```

INPUT

Enter the number of processes	4
Enter Burst Time for Process 0	6
Enter Burst Time for Process 1	8
Enter Burst Time for Process 2	7
Enter Burst Time for Process 3	3

OUTPUT

PROCESS	BURST	WAITING	TURNARO
	TIME	TIME	UND TIME
P3	3	0	3
P0	6	3	9
P2	7	9	16
P1	8	16	24
Average Waiting Time		7.000000	
Average Turnaround Time		13.000000	

C). ROUND ROBIN:

AIM: To simulate the CPU scheduling algorithm round-robin.

DESCRIPTION:

To aim is to calculate the average waiting time. There will be a time slice, each process should be executed within that time-slice and if not it will go to the waiting state so first check whether the burst time is less than the time-slice. If it is less than it assign the waiting time to the sum of the total times. If it is greater than the burst-time then subtract the time slot from the actual burst time and increment it by time-slot and the loop continues until all the processes are completed.

ALGORITHM:

- Step 1: Start the process
- Step 2: Accept the number of processes in the ready Queue and time quantum (or) time slice
- Step 3: For each process in the ready Q, assign the process id and accept the CPU burst time
- Step 4: Calculate the no. of time slices for each process where No. of time
- slice for process (n) = burst time process (n)/time slice
- Step 5: If the burst time is less than the time slice then the no. of time slices =1.
- Step 6: Consider the ready queue is a circular Q, calculate
- a) Waiting time for process (n) = waiting time of process(n-1)+ burst time of process(n-1) + the time difference in getting the CPU fromprocess(n-1)
- b) Turnaround time for process(n) = waiting time of process(n) + burst time of process(n)+ the time difference in getting CPU from process(n).

 Star 7: Colorlate
 - Step 7: Calculate
 - c) Average waiting time = Total waiting Time / Number of process
- d) Average Turnaround time = Total Turnaround Time / Number ofprocess Step 8: Stop the process

SOURCE CODE

```
#include<stdio.h>
main()
{
int
i,j,n,bu[10],wa[10],tat[10],t,ct[10],max; float
awt=0,att=0,temp=0;
clrscr();
printf("Enter the no of processes -- ");
scanf("%d",&n);
for(i=0;i<n;i++)
printf("\nEnter Burst Time for process %d -- ", i+1);
scanf("%d",&bu[i]);
ct[i]=bu[i];
}
printf("\nEnter the size of time slice -- ");
scanf("%d",&t);
max=bu[0];
for(i=1;i<n;i++)
if(max<bu[i])
max=bu[i];
for(j=0;j \le (max/t)+1;j++)
for(i=0;i \le n;i++)
if(bu[i]!=0)
if(bu[i] \le t) {
tat[i]=temp+bu[i];
temp=temp+bu[i];
bu[i]=0;
}
else {
bu[i]=bu[i]-t;
temp=temp+t;
}
for(i=0;i\le n;i++){
wa[i]=tat[i]-
ct[i]; att+=tat[i];
awt += wa[i];
printf("\nThe Average Turnaround time is -- %f",att/n);
printf("\nThe Average Waiting time is -- %f ",awt/n);
printf("\n\tPROCESS\t BURST TIME \t WAITING TIME\tTURNAROUND TIME\n");
for(i=0;i < n;i++)
printf("\t%d \t %d \t\t %d \n",i+1,ct[i],wa[i],tat[i]);
getch();}
```

INPUT:

Enter the no of processes -3

Enter Burst Time for process 1 – 24

Enter Burst Time for process 2 -- 3

Enter Burst Time for process 3 - 3

Enter the size of time slice -3

OUTPUT:

PROCESS	BURST TIME	WAITING TIME	TURNAROUNDTIM E
1	24	6	30
2	3	4	7
3	3	7	10

The Average Turnaround time is -15.666667 The

Average Waiting time is

5.666667

D). PRIORITY:

AIM: To write a c program to simulate the CPU scheduling priorityalgorithm.

DESCRIPTION:

To calculate the average waiting time in the priority algorithm, sort the burst times according to their priorities and then calculate the average waiting time of the processes. The waiting time of each process is obtained by summing up the burst times of all the previous processes.

ALGORITHM:

- Step 1: Start the process
- Step 2: Accept the number of processes in the ready Queue
- Step 3: For each process in the ready Q, assign the process id and accept the CPU burst time
- Step 4: Sort the ready queue according to the priority number.
- Step 5: Set the waiting of the first process as _0' and its burst time as its turnaround time
- Step 6: Arrange the processes based on process priority
- Step 7: For each process in the Ready Q calculate Step 8:

for each process in the Ready Q calculate

- a) Waiting time(n)= waiting time (n-1) + Burst time (n-1)
- b) Turnaround time (n)= waiting time(n)+Burst time(n)
- Step 9: Calculate
 - c) Average waiting time = Total waiting Time / Number of process
- d) Average Turnaround time = Total Turnaround Time / Number of process Print the results in an order.

Step10: Stop

SOURCE CODE:

```
#include<stdio.h>
main()
{
int p[20],bt[20],pri[20], wt[20],tat[20],i, k, n, temp; float wtavg,
tatavg;
clrscr();
printf("Enter the number of processes --- ");
scanf("%d",&n);
for(i=0;i<n;i++){
p[i] = i;
printf("Enter the Burst Time & Priority of Process %d --- ",i); scanf("%d
%d",&bt[i], &pri[i]);
for(i=0;i\le n;i++)
for(k=i+1;k \le n;k++)
if(pri[i] > pri[k]){
temp=p[i];
p[i]=p[k];
p[k]=temp;
temp=bt[i];
bt[i]=bt[k];
bt[k]=temp;
temp=pri[i];
pri[i]=pri[k];
pri[k]=temp;
wtavg = wt[0] = 0;
tatavg = tat[0] = bt[0];
for(i=1;i \le n;i++)
wt[i] = wt[i-1] + bt[i-1];
tat[i] = tat[i-1] + bt[i];
wtavg = wtavg + wt[i];
tatavg = tatavg + tat[i];
printf("\nPROCESS\t\tPRIORITY\tBURST TIME\tWAITING TIME\tTURNAROUND
TIME");
for(i=0;i<n;i++)
printf("\n%d \t\t %d \t\t fi],pri[i],br[i],wt[i],tat[i]);
printf("\nAverage Waiting Time is --- %f",wtavg/n); printf("\nAverage
Turnaround Time is --- %f",tatavg/n);
getch();}
```

INPUT

Enter the number of processes -- 5

Enter the Burst Time & Priority of Process 0 10	3
Enter the Burst Time & Priority of Process 1 1	1
Enter the Burst Time & Priority of Process 2 2	4
Enter the Burst Time & Priority of Process 3 1	5
Enter the Burst Time & Priority of Process 4 5	2

OUTPUT

PROCESS	PRIORITY	BURST TIME	WAITIN	TURNARO
			G TIME	UND TIME
1	1	I	0	1
4	2	5	1	6
0	3	10	6	16
2	4	2	16	18
3	5	1	18	19
Average Waiting	Time is	8.200000		
Average Turnarou	ınd Time is	12.000000		

VIVA QUESTIONS

- 1) Define the following
 - a) Turnaround time
- b) Waiting time
- c) Burst time
- d) Arrival time

- 2) What is meant by process scheduling?
- 3) What are the various states of process?
- 4) What is the difference between preemptive and non-preemptive scheduling
- 5) What is meant by time slice?
- 6) What is round robin scheduling?

EXPERIMENT.NO 2

AIM: To Write a C program to simulate producer-consumer problem using semaphores.

DESCRIPTION

Producer consumer problem is a synchronization problem. There is a fixed size buffer where the producer produces items and that is consumed by a consumer process. One solution to the producer-consumer problem uses shared memory. To allow producer and consumer processes to run concurrently, there must be available a buffer of items that can be filled by the producer and emptied by the consumer. This buffer will reside in a region of memory that is shared by the producer and consumer processes. The producer and consumer must be synchronized, so that the consumer does not try to consume an item that has not yet been produced.

PROGRAM

```
#include<stdio.
> void main()
1
       int buffer[10], bufsize, in, out, produce, consume,
       choice=0; in = 0;
       out = 0;
       bufsize = 10;
       while(choice !=3)
                            Produce
              printf("\n1.
                                          2. Consume \t3. Exit");
                                       \t
              printf("\nEnter your choice: ");
              scanf("%d",&choice);
              switch(choice) {
                      case 1: if((in+1)%bufsize=out)
                                     printf("\nBuffer is Full");
                             else
                                         printf("\nEnter the
                                         value: "); scanf("%d",
                                         &produce); buffer[in] =
                                         produce;
                                         in = (in+1)\%bufsize;
                             break;;;
                    case 2: if(in == out)
                                  printf("\nBuffer is Empty");
                                               else
                                               consume = buffer[out];
                                               printf("\nThe consumed
                                               value is %d", consume); out
                                               = (out+1)%bufsize;
                                               break;
```

OUTPUT

1. Produce 2. Consume 3. Exit Enter your choice: 2

Buffer is Empty

1. Produce 2. Consume 3.

Exit Enter your choice: 1

Enter the value: 100

1. Produce 2. Consume 3.

Exit Enter your choice: 2

The consumed value is 100

1. Produce 2. Consume 3.

Exit Enter your choice: 3

EXPERIMENT.NO 3

AIM: To Write a C program to simulate the concept of Dining-Philosophers problem.

DESCRIPTION

The dining-philosophers problem is considered a classic synchronization problem because it is an example of a large class of concurrency-control problems. It is a simple representation of the need to allocate several resources among several processes in a deadlock-free and starvation-free manner. Consider five philosophers who spend their lives thinking and eating. The philosophers share a circular table surrounded by five chairs, each belonging to one philosopher. In the center of the table is a bowl of rice, and the table is laid with five single chopsticks. When a philosopher thinks, she does not interact with her colleagues. From time to time, a philosopher gets hungry and tries to pick up the two chopsticks that are closest to her (the chopsticks that are between her and her left and right neighbors). A philosopher may pick up only one chopstick at a time. Obviously, she cam1ot pick up a chopstick that is already in the hand of a neighbor. When a hungry philosopher has both her chopsticks at the same time, she eats without releasing her chopsticks. When she is finished eating, she puts down both of her chopsticks and starts thinking again. The dining-philosophers problem may lead to a deadlock situation and hence some rules have to be framed to avoid the occurrence of deadlock.

PROGRAM

```
int tph, philname[20], status[20], howhung, hu[20], cho; main()
       int i; clrser();
       printf("\n\nDINING PHILOSOPHER PROBLEM");
       printf("\nEnter the total no. of philosophers: ");
       scanf("%d", &tph);
       for(i=0;i \le tph;i++)
              philname[i]=(i+1); status[i]=1;
       printf("How many are hungry:
        "); scanf("%d", &howhung);
        if(howhung==tph)
     printf("\n All are hungry..\nDead lock stage will occur");
     printf(\n"Exiting\n");
     else{
     for(i=0;i \leq howhung;i++)
     printf("Enterphilosopher%dposition:",(i+1));
     scanf("%d",&hu[i]);
     status[hu[i]]=2;
     }
```

```
do
               {
                      printf("1.One can eat at a time\t2.Two can eat at a time
                      \t3.Exit\nEnter your choice:");
                       scanf("%d", &cho);
                      switch(cho)
                       {
                       case 1: one();
                                                  break;
                    case 2: two();
                            break; case 3:
                    exit(0);
                                                                        default: printf("\nInvalid
                                          }
                                  }while(1);
                                                                        option..");
       }
}
one()
       int pos=0, x, i;
       printf("\nAllow one philosopher to eat at any time\n");
       for(i=0;i<howhung; i++, pos++)
       {
    printf("\nP %d is granted to eat", philname[hu[pos]]);
    for(x=pos;x<howhung;x++)
    printf("\nP %d is waiting", philname[hu[x]]);
                                                  }
}
two()
{
                                                  int i, j, s=0, t, r, x;
                                                  printf("\n
                                                  Allow two
                                                  philosophers
                                                  to eat at same
                                                  time\n");
                                                  for(i=0;i<how
                                                 hung;i++)
                                                  {
               for(j=i+1;j<howhung;j++)
               1
                      if(abs(hu[i]-hu[j]) \ge 1\&\& abs(hu[i]-hu[j])!=4)
                       {
                              printf("\n\ncombination %d \n", (s+1));
```

```
t=hu[i];
r=hu[j]; s++;
printf("\nP %d and P %d are granted to eat", philname[hu[i]],
     philname[hu[j]]);
```

```
for(x=0;x<howhung;x++)
                                  {
                                         if((hu[x]!=t)&&(hu[x]!=r))
                                         printf("\nP %d is waiting", philname[hu[x]]);
                                  }
                            }
                     }
              }
        }
                                                   DINING PHILOSOPHER PROBLEM
INPUT
                                                   E
                                                   h
                                                   e
                                                   t
                                                   0
                                                   t
                                                   a
                                                   1
                                                   n
                                                   o
                                                   0
                                                   f
                                                   p
                                                   h
                                                   i
                                                   1
                                                   o
                                                   S
                                                   0
                                                   p
                                                   h
                                                   e
```

r s

h : 5 u Η n 0 g r W y m a 3 n Enter philosopher 1 position: 2 y Enter philosopher 2 position: 4 a Enter philosopher 3 position: 5 r e

OUTPUT

1. One can eat at a time 2. Two can eat at a time 3. Exit Enter your choice: 1

Allow one philosopher to eat at any time P 3 is granted to eat

P 3 is waiting

P 5 is waiting

P 0 is waiting

P 5 is granted to eat

P 5 is waiting

P 0 is waiting

P 0 is granted to eat

P 0 is waiting

1.One can eat at a time

2.Two can eat at a time 3.Exit

Enter your choice: 2

Allow two philosophers to eat at same time combination 1
P 3 and P 5 are granted to eat
P 0 is waiting

combination 2

P 3 and P 0 are granted to eat

P 5 is waiting

combination 3

P 5 and P 0 are granted to eat

P 3 is waiting

1.One can eat at a time

2.Two can

eat at a time

3.Exit Enter your choice: 3

EXPERIMENT.NO 4 MEMORY MANAGEMENT

A). MEMORY MANAGEMENT WITH FIXED PARTITIONING TECHNIQUE (MFT)

AIM: To implement and simulate the MFT algorithm.

DESCRIPTION:

In this the memory is divided in two parts and process is fit into it. The process which is best suited will be placed in the particular memory where it suits. In MFT, the memory is partitioned into fixed size partitions and each job is assigned to a partition. The memory assigned to a partition does not change. In MVT, each job gets just the amount of memory it needs. That is, the partitioning of memory is dynamic and changes as jobs enter and leave the system. MVT is a more ``efficient" user of resources. MFT suffers with the problem of internal fragmentation and MVT suffers with external fragmentation.

ALGORITHM:

Step 8: Print the results

```
Step1: Start the process.
Step2: Declarevariables.
Step3: Enter total memory size ms.
Step4: Allocate memory for os.
Ms=ms-os
Step5: Read the no partition to be divided n Partition size=ms/n.
Step6: Read the process no and process size.
Step 7: If process size is less than partition size allot alse blocke the process. While allocating update memory wastage-external fragmentation.
if(pn[i]==pn[j])f=1;
if(f==0){ if(ps[i]<=siz) {
    extft=extft+size-
    ps[i];avail[i]=1; count++;
}
```

SOURCE CODE:

```
#include<stdio.h>
#include < conio.h >
main()
1
int
      ms,
              bs,
                    nob,
ef,n, mp[10],tif=0; int i,p=0;
clrscr();
printf("Enter the total memory available (in Bytes) -- ");
scanf("%d",&ms);
printf("Enter the block size (in Bytes) -- ");
scanf("%d", &bs);
nob=ms/bs;
ef=ms - nob*bs;
printf("\nEnter the number of processes -- ");
scanf("%d",&n);
for(i=0;i\leq n;i++)
printf("Enter memory required for process %d (in Bytes)-- ",i+1);
scanf("%d",&mp[i]);
}
printf("\nNo.
                                Blocks
                                                                            memory--%d",nob);
                     of
                                                available
                                                                 in
printf("\n\nPROCESS\tMEMORYREQUIRED\tALLOCATED\tINTERNAL
FRAGMENTATION");
for(i=0;i<n && p<nob;i++)
printf("\n %d\t\t%d",i+1,mp[i]);
if(mp[i] > bs)
printf("\t\tNO\t\t---");
else
1
printf("\t\tYES\t%d",bs-mp[i]);
tif = tif + bs-mp[i];
p++;
}
if(i \le n)
printf("\nMemory is Full, Remaining Processes cannot be accomodated");
printf("\n\nTotal Internal Fragmentation is %d",tif);
printf("\nTotal External Fragmentation is %d",ef);
getch();
```

INPUT

Enter the total memory available (in Bytes)	1000
Enter the block size (in Bytes) 300	
Enter the number of processes -5	
Enter memory required for process 1 (in Bytes)	275
Enter memory required for process 2 (in Bytes)	400
Enter memory required for process 3 (in Bytes)	290
Enter memory required for process 4 (in Bytes)	293
Enter memory required for process 5 (in Bytes)	100
No. of Blocks available in memory 3	

OUTPUT

PROCESS		ALLOCAT	INTERNAL
	MEMORY REQUIRED	ED	FRAGMENTATIO
			N
1	275	YES	25
2	400	NO	
3	290	YES	10
4	293	YES	7

Memory is Full, Remaining Processes cannot be accommodated Total Internal Fragmentation is 42

Total External Fragmentation is 100

B) MEMORY VARIABLE PARTIONING TYPE (MVT)

AIM: To write a program to simulate the MVT algorithm

ALGORITHM:

Step1: start the process. Step2: Declare variables.

Step3: Enter total memory size ms.

Step4: Allocate memory for os.

Ms=ms-os

Step 5: Read the no partition to be divided n Partition size=ms/n.

Step6: Read the process no and process size.

Step 7: If process size is less than partition size allot alse blocke the process. While allocating update memory wastage-external fragmentation.

```
if(pn[i]==pn[j]) f=1;
if(f==0){ if(ps[i]<=size)
{
  extft=extft+size-
ps[i];avail[i]=1; count++;
}
}</pre>
```

Step 8: Print the results

Step 9: Stop the process.

SOURCE CODE:

```
#include<stdio.h>
#include<conio.h>
main()
{
int
ms,mp[10],i,temp,n=0;
char ch = 'y'; clrscr();
printf("\nEnter the total memory available (in Bytes)-- ");
scanf("%d",&ms);
temp=ms;
for(i=0;ch=='y';i++,n++)
1
printf("\nEnter memory required for process %d (in Bytes) -- ",i+1);
scanf("%d",&mp[i]);
if(mp[i] \le temp)
{
printf("\nMemory is allocated for Process %d ",i+1);
temp = temp - mp[i];
}
else
printf("\nMemory is Full"); break;
printf("\nDo you want to continue(y/n) -- ");
 scanf(" %c", &ch);
printf("\n\nTotal
                   Memory Available
                                                %d",
                                                        ms);
ALLOCATED
for(i=0;i\leq n;i++)
printf("\n \t\%d\t\t\%d",i+1,mp[i]);
printf("\n\nTotal
                  Memory
                               Allocated
                                                 %d",ms-temp);
                                            is
printf("\nTotal External Fragmentation is %d",temp);
getch();
```

OUTPUT:

Enter the total memory available (in Bytes) – 1000
Enter memory required for process 1 (in Bytes) – 400
Memory is allocated for Process 1
Do you want to continue(y/n) -- y
Enter memory required for process 2 (in Bytes) -- 275
Memory is allocated for Process 2
Do you want to continue(y/n) – y
Enter memory required for process 3 (in Bytes) – 550

Memory is Full

Total Memory Available – 1000

PROCESS MEMORY ALLOCATED
400
2 275

Total Memory Allocated is 675
Total External Fragmentation is 325

VIVA QUESTIONS

- 1) What is MFT?
- 2) What is MVT?
- 3) What is the difference between MVT and MFT?
- 4) What is meant by fragmentation?
- 5) Give the difference between internal and external fragmentation

EXPERIMENT.NO 5 MEMORY ALLOCATION TECHNIQUES

AIM: To Write a C program to simulate the following contiguous memory allocation techniques a) Worst-fit b) Best-fit c) First-fit

DESCRIPTION

One of the simplest methods for memory allocation is to divide memory into several fixed-sized partitions. Each partition may contain exactly one process. In this multiple-partition method, when a partition is free, a process is selected from the input queue and is loaded into the free partition. When the process terminates, the partition becomes available for another process. The operating system keeps a table indicating which parts of memory are available and which are occupied. Finally, when a process arrives and needs memory, a memory section large enough for this process is provided. When it is time to load or swap a process into main memory, and if there is more than one free block of memory of sufficient size, then the operating system must decide which free block to allocate. Best-fit strategy chooses the block that is closest in size to the request. First-fit chooses the first available block that is large enough. Worst-fit chooses the largest available block.

PROGRAM

WORST-FIT

```
#include<stdio.h>
#include<conjo.h>
#define max 25
void main()
1
       int
       frag[max],b[max],f[max],i,j,nb,nf,t
       emp; static int bf[max],ff[max];
       clrscr();
       printf("\n\tMemory Management Scheme - First
       Fit"); printf("\nEnter the number of blocks:");
       scanf("%d",&nb);
       printf("Enter the number of files:");
       scanf("%d",&nf);
       printf("\nEnter the size of the
       blocks:-\n"); for(i=1;i<=nb;i++)
              printf("Block %d:",i);
              scanf("%d",&b[i]);
       printf("Enter
                      the
                             size
                                   of
                                        the
                                              files
                                                     :-\n");
       for(i=1;i \le nf;i++)
       {
              printf("File %d:",i);
              scanf("%d",&f[i]);
```

```
}
        for(i=1;i \le nf;i++)
        {
               for(j=1;j<=nb;j++)
                       if(bf[j]!=1)
                        {
                               temp=b[j]-f[i]
                               ; if(temp\geq = 0)
                                       ff[i] = j
                                       break;
                             }
                     }
             frag[i]=temp
             ; bf[ff[i]]=1;
        }
                                                  printf("\nFile no:\tFile size
                                                  :\tBlock no:\tBlock size:\tFra
                                                  gement"); for(i=1;i \le nf;i++)
                                                  printf("\n%d\t\t%d\t\t%d\t\t%d
                                                  \t\t%d",i,f[i],ff[i],b[ff[i]],frag[i]
                                                  ); getch();
Enter the number of blocks: 3
Enter the number of files: 2
Enter the size of the blocks:-
Block 1:5
Block 2: 2
Block 3: 7
Enter the size of the files:-
File 1: 1
File 2: 4
```

INPUT

OUTPUT File No

File Size

1

4

Block No

1

3

Block Size

5

7

Fragment

4

3

```
BEST-FIT
```

```
#include<stdio.h>
#include<conio.h>
#define max 25
void main()
{
        int frag[max],b[max],f[max],i,j,nb,nf,temp,lowest=10000;
        static int bf[max],ff[max];
        clrscr();
       printf("\nEnter the number of blocks:");
        scanf("%d",&nb);
        printf("Enter the number of files:");
        scanf("%d",&nf);
        printf("\nEnter the size of the
        blocks:-\n"); for(i=1;i \le nb;i++)
     printf("Block %d:",i);
     scanf("%d",&b[i]);
       printf("Enter the size of the
                                                     :-\n");
                                             files
       for(i=1;i<=nf;i++)
       {
               printf("File %d:",i);
               scanf("%d",&f[i]);
       }
        for(i=1;i \le nf;i++)
                for(j=1;j\leq nb;j++)
                      if(bf[j]!=1)
                               temp=b[j]-f[i]
                               ; if(temp\geq = 0)
                                      if(lowest>temp)
                                      ff[i]=j;
                                       lowest=temp;
                    }}
                frag[i]=lowest; bf[ff[i]]=1; lowest=10000;
        printf("\nFile No\tFile Size \tBlock
        No\tBlock Size\tFragment"); for(i=1;i<=nf &&
        ff[i]!=0;i++)
               printf("\n%d\t\t%d\t\t%d\t\t%d\t\t%d\t\t%d",i,f[i],ff[i],b[ff[i]],frag[i]);
        getch();
}
```

INPUT

```
Enter the number of blocks: 3
Enter the number of files: 2

Enter the size of the blocks:-
Block 1: 5
Block 2: 2
Block 3: 7

Enter the size of the files:-
File 1: 1
File 2: 4

OUTPUT
File No File Size Block No
```

FIRST-FIT

1

4

}

1

2

```
#include<stdio.h>
#include<conio.h>
#define max 25
void main()
{
        int
        frag[max],b[max],f[max],i,j,nb,nf,temp,highes
        t=0; static int bf[max],ff[max];
        clrscr();
        printf("\n\tMemory Management Scheme - Worst Fit");
        printf("\nEnter the number of blocks:");
        scanf("%d",&nb);
        printf("Enter the number of files:");
        scanf("%d",&nf);
        printf("\nEnter the size of the blocks:-\n");
        for(i=1;i \le nb;i++)
               printf("Block %d:",i);
               scanf("%d",&b[i]);
        printf("Enter the size of the files :-\n");
        for(i=1;i \le nf;i++)
        {
               printf("File %d:",i);
               scanf("%d",&f[i]);
```

2

1

Block Size

Fragment

1

1

```
for(i=1;i \le nf;i++)
               for(j=1;j\le nb;j++)
                       if(bf[j]!=1) //if bf[j] is not allocated
                       {
                              temp=b[j]-f[i]
                              ; if(temp\geq =0)
                                      if(highest<temp)
                       }
                frag[i]=highest; bf[ff[i]]=1; highest=0;
        ff[i]=j; highest=temp;
     }
        printf("\nFile no:\tFile size:\tBlock no:\tBlock size:\tFragement");
        for(i=1;i \le nf;i++)
               printf("\n\%d\t\t\%d\t\t\%d\t\t\%d",i,f[i],ff[i],b[ff[i]],frag[i]);
        getch();
}
Enter the number of blocks: 3
Enter the number of files: 2
Enter the size of the blocks:-
Block 1:5
Block 2: 2
Block 3: 7
Enter the size of the files:-
File 1: 1
File 2: 4
OUTPUT
File No
               File Size
                               Block No
                                              Block Size
                                                              Fragment
               1
                               3
                                               7
                                                              6
```

5

1

INPUT

4

1

EXPERIMENT NO.6

PAGE REPLACEMENT ALGORITHMS

AIM: To implement FIFO page replacement technique.

a) FIFO b) LRU c) OPTIMAL

DESCRIPTION:

Page replacement algorithms are an important part of virtual memory management and it helps the OS to decide which memory page can be moved out making space for the currently needed page. However, the ultimate objective of all page replacement algorithms is to reduce the number of page faults.

FIFO-This is the simplest page replacement algorithm. In this algorithm, the operating system keeps track of all pages in the memory in a queue, the oldest page is in the front of the queue. When a page needs to be replaced page in the front of the queue is selected for removal.

LRU-In this algorithm page will be replaced which is least recently used

OPTIMAL- In this algorithm, pages are replaced which would not be used for the longest duration of time in the future. This algorithm will give us less page fault when compared to other page replacement algorithms.

ALGORITHM:

- 1. Start the process
- 2. Read number of pages n
- 3. Read number of pages no
- 4. Read page numbers into an array a[i]
- 5. Initialize avail[i]=0 .to check page hit
 - 6. Replace the page with circular queue, while re-placing check page availability in the frame Place avail[i]=1 if page is placed in the frame Count page faults
- 7. Print the results.
- 8. Stop the process.

A) FIRST IN FIRST OUT SOURCE CODE:

```
#include<stdio.h>
#include < conio.h > int fr[3];
void main()
void display();
int i,j,page[12]=\{2,3,2,1,5,2,4,5,3,2,5,2\};
flag1=0,flag2=0,pf=0,frsize=3,top=0;
clrscr();
for(i=0;i<3;i++)
fr[i]=-1;
for(j=0;j<12;j++)
flag1=0; flag2=0; for(i=0;i<12;i++)
if(fr[i]==page[j])
flag1=1; flag2=1; break;
if(flag1==0)
for(i=0;i<frsize;i++)
if(fr[i]==-1)
fr[i]=page[j]; flag2=1; break;
if(flag2==0)
fr[top]=page[j];
top++;
pf++;
if(top>=frsize)
top=0;
display();
```

```
\label{eq:printf} $$ printf("Number of page faults : %d ",pf+frsize); $$ getch(); $$ void display() $$ {$ int i; $ printf("\n"); $$ for(i=0;i<3;i++) $$ printf("%d\t",fr[i]); $$ $$ $$ $$
```

2 -1 -1

2 3 -1

2 3 -1

2 3 1

5 3 1

5 2 1

5 2 4

5 2 4

3 2 4

3 2 4

3 5 4

3 5 2

Number of page faults: 9

B) LEAST RECENTLY USED

AIM: To implement LRU page replacement technique.

ALGORITHM:

- 1. Start the process
- 2. Declare the size
- 3. Get the number of pages to be inserted
- 4. Get the value
- 5. Declare counter and stack
- 6. Select the least recently used page by counter value
- 7. Stack them according the selection.
- 8. Display the values
- 9. Stop the process

```
#include<stdio.h>
#include<conio.h>
int fr[3];
void main()
{
void display();
int p[12] = \{2,3,2,1,5,2,4,5,3,2,5,2\},i,j,fs[3];
int index,k,l,flag1=0,flag2=0,pf=0,frsize=3;
elrser();
for(i=0;i<3;i++)
fr[i]=-1;
for(j=0;j<12;j++)
flag1=0,flag2=0;
for(i=0;i<3;i++)
if(fr[i]==p[j])
flag1=1;
flag2=1; break;
if(flag1==0)
```

```
for(i=0;i<3;i++)
if(fr[i]==-1)
fr[i]=p[j];
flag2=1; break;
if(flag2==0)
for(i=0;i<3;i++)
fs[i]=0;
for(k=j-1,l=1;l<=frsize-1;l++,k--)
for(i=0;i<3;i++)
if(fr[i]==p[k]) fs[i]=1;
}}
for(i=0;i<3;i++)
if(fs[i]==0)
index=i;
}
fr[index]=p[j];
pf++;
display();
printf("\n no of page faults :%d",pf+frsize);
getch();
}
void display()
int i; printf("\n");
for(i=0;i<3;i++)
printf("\t%d",fr[i]);
```

- 2 -1 -1
- 2 3 -1
- 2 3 -1
- 2 3 1
- 2 5 1
- 2 5 1
- 2 5 4
- 2 5 4
- 3 5 4
- 3 5 2
- 3 5 2
- 3 5 2

No of page faults: 7

C) OPTIMAL

AIM: To implement optimal page replacement technique.

ALGORTHIM:

- 1. Start Program
- 2. Read Number Of Pages And

Frames 3.Read Each Page Value

- 4. Search For Page In The Frames
- 5.If Not Available Allocate Free Frame
 - 6. If No Frames Is Free Repalce The Page With The Page That Is Leastly Used7.Print Page Number Of Page Faults8.Stop process.

```
Program to
                   simulate
                              optimal page replacement
*/#include<stdio.h>
#include<conio.h>
int fr[3], n, m;
void
display();
void main()
int i,j,page[20],fs[10];
max,found=0,lg[3],index,k,l,flag1=0,flag2=0,pf=0
; float pr;
clrscr();
printf("Enter length of the
                                reference string:
"); scanf("%d",&n);
printf("Enter the
                     reference
                                 string:
                                           ");
for(i=0;i\le n;i++)
scanf("%d",&page[i]);
printf("Enter no of frames:
                                  ");
scanf("%d",&m);
for(i=0;i<m;i++)
fr[i]=-1; pf=m;
```

```
for(j=0;j<n;j++)
flag1=0;
flag2=0;
for(i=0;i\le m;i++)
if(fr[i]==page[j])
flag1=1; flag2=1;
break;
}
if(flag1==0)
for(i=0;i\le m;i++)
if(fr[i]==-1)
fr[i]=page[j]; flag2=1;
break;
}
if(flag2==0)
for(i=0;i<m;i++)
lg[i]=0;
for(i=0;i<m;i++)
{
for(k=j+1;k<=n;k++)
if(fr[i]==page[k])
lg[i]=k-j;
break;
found=0;
for(i=0;i<m;i++)
if(lg[i]==0)
index=i;
found = 1;
```

```
break;
}
}
if(found==0)
max=lg[0]; index=0;
for(i=0;i<m;i++)
{
if(max{<}lg[i])\\
max=lg[i];
index=i;
fr[index]=page[j];
pf++;
}
display();
printf("Number of page faults : %d\n",
pf); pr=(float)pf/n*100;
printf("Page fault rate = %f \n", pr); getch();
void display()
{
int i;
for(i=0;i<m;i++)
printf("%d\t",fr[i]);
printf("\n");
}
```

Enter length of the reference string: 12

Enter the reference string: 1 2 3 4 1 2 5 1 2 3 4 5

Enter no of frames: 3

1 -1 -1

12-1

123

124

124

124

125

1 2 5

125

3 2 5

425

425

Number of page faults: 7 Page fault rate = 58.333332

VIVA QUESTIONS

- 1) What is meant by page fault?
- 2) What is meant by paging?
- 3) What is page hit and page fault rate?
- 4) List the various page replacement algorithm
- 5) Which one is the best replacement algorithm?

EXPERIMENT NO. 7

FILE ORGANIZATION TECHNIQUES

A) SINGLE LEVEL DIRECTORY:

AIM: Program to simulate Single level directory file organization technique.

DESCRIPTION:

The directory structure is the organization of files into a hierarchy of folders. In a single-level directory system, all the files are placed in one directory. There is a root directory which has all files. It has a simple architecture and there are no sub directories. Advantage of single level directory system is that it is easy to find a file in the directory.

```
#include<stdio.h>
struct
1
char
dname[10],fname[10][10]; int
fent;
}dir;
void main()
int i,ch; char
f[30]; clrscr();
dir.fent = 0;
                             directory --
printf("\nEnter name of
"); scanf("%s", dir.dname);
while(1)
printf("\n\n1. Create File\t2. Delete File\t3. Search File \n
4. Display Files\t5. Exit\nEnter your choice -- ");
scanf("%d",&ch);
switch(ch)
case 1: printf("\nEnter the name of the file -- ");
scanf("%s",dir.fname[dir.fcnt]);
dir.fcnt++; break;
case 2: printf("\nEnter the name of the file -- ");
scanf("%s",f);
for(i=0;i<dir.fcnt;i++)
if(strcmp(f, dir.fname[i])==0)
printf("File %s is deleted ",f); strcpy(dir.fname[i],dir.fname[dir.fcnt-1]); break;
```

```
if(i=dir.fcnt)
printf("File %s not
                               else
found",f);
                                                                      k
                                           printf("\nEnter the name of the file -- ");
                                case 3:
                                     scanf("%s",f);
                                     for(i=0;i<dir.fcnt;i++)
                                     if(strcmp(f, dir.fname[i])==0)
                                     printf("File %s is found ", f);
                                     break;
                                     if(i=dir.fcnt)
                                     printf("File %s not found",f);
                                     break;
                                case 4:
                                          if(dir.fcnt=0)
                                     printf("\nDirectory Empty");
                                     else.
                                     printf("\nThe Files are -- ");
                                     for(i=0;i<dir.fcnt;i++)
                                     printf("\t%s",dir.fname[i]);
                                     break;
                                                    getch();}
```

```
default: exit(0);
}
```

Enter name of directory -- CSE

- 1. Create File 2. Delete File 3. Search File
- 4. Display Files 5. Exit Enter your choice 1

Enter the name of the file -- A

- 1. Create File 2. Delete File 3. Search File
- 4. Display Files 5. Exit Enter your choice 1

Enter the name of the file -- B

- 1. Create File 2. Delete File 3. Search File
- 4. Display Files 5. Exit Enter your choice 1

Enter the name of the file -- C

- 1. Create File 2. Delete File 3. Search File
- 4. Display Files 5. Exit Enter your choice 4

The Files are -- ABC

- 1. Create File 2. Delete File 3. Search File
- 4. Display Files 5. Exit Enter your choice 3

Enter the name of the file – ABC File

ABC not found

- 1. Create File 2. Delete File 3. Search File
- 4. Display Files 5. Exit Enter your choice 2

Enter the name of the file - B

File B is deleted

- 1. Create File 2. Delete File 3. Search File
- 4. Display Files 5. Exit Enter your choice 5

B) TWO LEVEL DIRECTORY

AIM: Program to simulate two level file organization technique

Description:

In the two-level directory system, each user has own user file directory (UFD). The system maintains a master block that has one entry for each user. This master block contains the addresses of the directory of the users. When a user job starts or a user logs in, the system's master file directory (MFD) is searched. When a user refers to a particular file, only his own UFD is searched.

```
#include<stdio.h>
struct
1
       char
       dname[10],fname[10][10]; int
}dir[10];
void main()
             i,ch,dent,k;
       char f[30], d[30];
       clrser(); dent=0;
       while(1)
       {
               printf("\n\n1. Create Directory\t2. Create File\t3. Delete File");
               printf("\n4. Search File\t\t5. Display\t6. Exit\t Enter your choice --");
               scanf("%d",&ch);
               switch(ch)
                       case 1: printf("\nEnter name of directory -- ");
                               scanf("%s",
                                                     dir[dent].dname);
                               dir[dent].fent=0;
                               dent++;
                               printf("Directory created"); break;
                       case 2: printf("\nEnter name of the directory -- ");
                              scanf("%s",d);
                               for(i=0;i < dent;i++)
                                      if(strcmp(d,dir[i].dname)==0)
                                      1
                         printf("Enter name
                                                                 file
                                                 of
                         "); scanf("%s",dir[i].fname[dir[i].fcnt]);
```

```
dir[i].fcnt++;
                   printf("File created");
               if(i==dent)
                  printf("Directory %s not found",d);
       case 3: printf("\nEnter name of the directory --
               scanf("%s",d);
               for(i=0;i < dent;i++)
               for(i=0;i<dent;i++)
               {
               if(strcmp(d,dir[i].dname)==0)
                       printf("Enter name of the file -- ");
                       scanf("%s",f);
                       for(k=0;k\leq dir[i].fcnt;k++)
                       if(stremp(f, dir[i].fname[k])==0)
                       printf("File %s is deleted ",f);
                       dir[i].fcnt--;
                       strcpy(dir[i].fname[k],dir[i].fname[dir[i].fcnt]);
                       goto jmp;
               printf("File %s not found",f); goto imp;
       printf("Directory %s not found",d);
       jmp: break;
case 4: printf("\nEnter name of the directory -- ");
       scanf("%s",d);
       for(i=0;i<dent;i++)
               if(strcmp(d,dir[i].dname)==0)
                       printf("Enter the name of the file -- ");
                       scanf("%s",f);
                       for(k=0;k\leq dir[i].fent;k++)
                       {
                               if(stremp(f, dir[i].fname[k])==0)
                               printf("File %s is found ",f); goto jmp1;
                       }
       printf("
```

File %s

not found",f



```
printf("Directory %s not found",d); jmp1: break;
 case 5: if(dcnt==0)
printf("\nNo Directory's ");
        else
         {
                printf("\nDirectory\tFiles");
                for(i=0;i<dent;i++)
                {
                                 break;
                     default:exit(0);
             }
                                                    n
                                                    %
  getch();
                                                    S
                                                    d
                                                    i
                                                    ]
                                                    d
                                                    n
                                                    a
                                                    m
                                                    e
                                                    )
                                                    0
                                                    r
                                                    k
                                                    0
```

}

) ; ; k < d i r [i] . f c n t ; k +) p r i n t f (" t % S , d i [i] . f n a m e [k]

- 1. Create Directory 2. Create File 3. Delete File
- 4. Search File 5. Display 6. Exit

Enter your choice -- 1

Enter name of directory -- DIR1 Directory created

- 1. Create Directory 2. Create File 3. Delete File
- 4. Search File 5. Display 6. Exit Enter your choice -- 1

Enter name of directory -- DIR2 Directory created

- 1. Create Directory 2. Create File 3. Delete File
- 4. Search File 5. Display 6. Exit Enter your choice -- 2

Enter name of the directory – DIR1

Enter name of the file -- A1

File created

- 1. Create Directory 2. Create File 3. Delete File
- 4. Search File 5. Display 6. Exit

Enter your choice -- 2

Enter name of the directory - DIR1

Enter name of the file -- A2

File created

- 1. Create Directory 2. Create File 3. Delete File
- 4. Search File 5. Display 6.

Exit Enter your choice – 6

VIVA QUESTIONS

- 1. Define directory?
- 2. List the different types of directory structures?
- 3. What is the advantage of hierarchical directory structure?
- 4. Which of the directory structures is efficient? Why?
- 5. What is acyclic graph directory?

EXPERIMENT.NO.8 FILE ALLOCATION STRATEGIES

A) SEQUENTIAL:

AIM: To write a C program for implementing sequential file allocation method

DESCRIPTION:

The most common form of file structure is the sequential file in this type of file, a fixed format is used for records. All records (of the system) have the same length, consisting of the same number of fixed length fields in a particular order because the length and position of each field are known, only the values of fields need to be stored, the field name and length for each field are attributes of the file structure.

ALGORITHM:

```
Step 1: Start the program.
```

Step 2: Get the number of files.

Step 3: Get the memory requirement of each file.

Step 4: Allocate the required locations to each in sequential order a).

Randomly select a location from available location s1= random(100);

a) Check whether the required locations are free from the selected location.

```
if(b[s1].flag==0){
for
(j=s1;j<s1+p[i];j++){
    if((b[j].flag)==0)count++;
      }
    if(count==p[i]) break;
}
b) Allocate and set flag=1 to the allocated locations. for(s=s1;s<(s1+p[i]);s++)
{
    k[i][j]=s; j=j+1; b[s].bno=s;
    b[s].flag=1;
}</pre>
```

Step 5: Print the results file no, length, Blocks allocated. Step 6: Stop the program

```
#include<stdio.h>
main()
{
int f[50],i,st,j,len,c,k;
clrscr();
for(i=0;i<50;i++)
f[i]=0;
X:
printf("\n Enter the starting block & length of file");
scanf("%d%d",&st,&len);
for(j=st;j\leq(st+len);j++)
if(f[j]==0)
{
f[j]=1
printf("\n%d->%d",j,f[j]);
else
printf("Block already allocated");
break;
if(j==(st+len))
printf("\n the file is allocated to disk");
printf("\n if u want to enter more files?(y-1/n-0)");
scanf("%d",&c);
if(c==1)
goto X;
else
exit();
getch();
```

Enter the starting block & length of file 4

10 4->1

5->1

6->1

7->1

8->1

9->1

10->1

11->1

12->1

13->1

The file is allocated to disk.

B) INDEXED:

AIM: To implement allocation method using chained method

DESCRIPTION:

In the chained method file allocation table contains a field which points to starting block of memory. From it for each bloc a pointer is kept to next successive block. Hence, there is no external fragmentation.

ALGORITHM:

- Step 1: Start the program.
- Step 2: Get the number of files.
- Step 3: Get the memory requirement of each file.
- Step 4: Allocate the required locations by selecting a location randomly q= random(100);
 - a) Check whether the selected location is free.
 - b) If the location is free allocate and set flag=1 to the allocated locations.

```
q=random(100);
{
if(b[q].flag==0)
b[q].flag=1;
b[q].fno=j;
r[i][j]=q;
Step 5: Print the results file no, length ,Blocks allocated.
Step 6: Stop the program
```

```
#include<stdio.h>
int
f[50],i,k,j,inde[50],n,c,count=0,p;
main()
{
clrscr();
for(i=0;i<50;i++
) f[i]=0;
x: printf("enter index block\t");
scanf("%d",&p);
if(f[p]==0)
{
f[p]=1;
printf("enter no of files on index\t");
scanf("%d",&n);
}
else
{
printf("Block already allocated\n");
goto x;
}
for(i=0;i\leq n;i++)
scanf("%d",&inde[i]);
for(i=0;i<n;i++)
if(f[inde[i]]==1)
printf("Block already allocated");
goto x;
}
for(j=0;j \le n;j++)
f[inde[j]]=1;
printf("\n
              allocated");
printf("\n file indexed");
for(k=0;k\leq n;k++)
printf("\n %d->%d:%d",p,inde[k],f[inde[k]]);
printf(" Enter 1 to enter more files and 0 to exit\t");
seanf("%d",&c);
if(c==1)
goto
x; else
exit();
getch();
}
```

OUTPUT: enter index block 9
Enter no of files on index 3 1
2 3
Allocated
File indexed
9->1:1
9->2;1
9->3:1 enter 1 to enter more files and 0 to exit

C) LINKED:

AIM: To implement linked file allocation technique.

DESCRIPTION:

In the chained method file allocation table contains a field which points to starting block of memory. From it for each bloc a pointer is kept to next successive block. Hence, there is no external fragmentation

ALGORTHIM:

```
Step 1: Start the program.
Step 2: Get the number of files.
```

Step 3: Get the memory requirement of each file.

Step 4: Allocate the required locations by selecting a location randomly q=random(100);

- a) Check whether the selected location is free.
- b) If the location is free allocate and set flag=1 to the allocated locations. While allocating next location address to attach it to previous location

```
#include<stdio.h>
main()
1
int f[50],p,i,j,k,a,st,len,n,c;
clrscr();
for(i=0;i<50;i++) f[i]=0;
printf("Enter how many blocks that are already
allocated"); scanf("%d",&p);
printf("\nEnter the blocks no.s that are already allocated");
for(i=0;i<p;i++)
{
scanf("%d",&a);
f[a]=1;
}
X:
printf("Enter the starting index block
& length"); scanf("%d%d",&st,&len); k=len;
for(j=st;j<(k+st);j++)
{
if(f[j]==0)
\{f[j]=1;
printf("\n\%d->\%d",j,f[j]);
}
else
printf("\n %d->file is already
allocated",j);
k++;
}
printf("\n If u want to enter one
more file? (yes-1/no-0)");
scanf("%d",&c);
if(c=1)
goto
Χ;
else
exit();
getch();}
```

Enter how many blocks that are already allocated 3 Enter the blocks no.s that are already allocated 4.7 Enter the starting index block & length 3.7.9

3->1

4->1 file is already allocated 5->1

6->1

7->1 file is already allocated 8->1 9->1file is already allocated 10->1 11->1

12->1

VIVA QUESTIONS

- 1) List the various types of files
- 2) What are the various file allocation strategies?
- 3) What is linked allocation?
- 4) What are the advantages of linked allocation?
- 5) What are the disadvantages of sequential allocation methods?

EXPERIMENT.NO 9 DEAD LOCK AVOIDANCE

AIM: To Simulate bankers algorithm for Dead Lock Avoidance (Banker's Algorithm)

DESCRIPTION:

Deadlock is a situation where in two or more competing actions are waiting f or the other to finish, and thus neither ever does. When a new process enters a system, it must declare the maximum number of instances of each resource type it needed. This number may exceed the total number of resources in the system. When the user request a set of resources, the system must determine whether the allocation of each resources will leave the system in safe state. If it will the resources are allocation; otherwise the process must wait until some other process release the resources.

Data structures

- •n-Number of process, m-number of resource types.
- •Available: Available[j]=k, k instance of resource type Rj is available. Max:
- If •max[i, j]=k, Pi may request at most k instances resource Rj.
 - *Allocation: If Allocation [i, j]=k, Pi allocated to k instances of resource Rj Need: If
- Need[I, j]=k, Pi may need k more instances of resource type Rj, Need[I, j]=Max[I, j]-Allocation[I, j];

Safety Algorithm

- 1. Work and Finish be the vector of length m and n respectively, Work=Available and Finish[i] =False.
- 2. Find an i such that both
 - Finish[i] =False
 - •Need<=Work If

no such I exists go to step 4.

- 3. work= work + Allocation, Finish[i] =True;
 - 4. if Finish[1]=True for all I, then the system is in safe state.

Resource request algorithm

Let Request i be request vector for the process Pi, If request i=[j]=k, then process Pi wants k instances of resource type Rj.

- 1. if Request <= Need I go to step 2. Otherwise raise an error condition.
- 2. if Request = Available go to step 3. Otherwise Pi must since the resources are available.
- 3. Have the system pretend to have allocated the requested resources to process Pi by modifying the state as follows;

Available=Available-Request I;

Allocation I=Allocation +Request I;

Need i=Need i- Request I;

If the resulting resource allocation state is safe, the transaction is completed and process Pi is allocated its resources. However if the state is unsafe, the Pi must wait for Request i and the old resource-allocation state is restored.

ALGORITHM:

- 1. Start the program.
- 2. Get the values of resources and processes.
- 3. Get the avail value.
- 4. After allocation find the need value.
- 5. Check whether its possible to allocate.
- 6. If it is possible then the system is in safe state.
- 7. Else system is not in safety state.
- 8. If the new request comes then check that the system is in safety.
- 9. or not if we allow the request.
- 10. stop the program.
- 11. end

```
#include<stdio.h>
#include<conio.h>
#include<string.h>
void main()
{
int alloc[10][10],max[10][10];
int avail[10],work[10],total[10];
int i,j,k,n,need[10][10];
int m:
int
       count=0, c=0;
char
finish[10]; clrscr();
printf("Enter
                the
                       no.
                             of
                                   processes
                                                 and
resources:"); scanf("%d%d",&n,&m);
for(i=0;i \le n;i++)
finish[i]='n';
                   the
                            claim
printf("Enter
                                       matrix:\n");
for(i=0;i\leq n;i++)
for(j=0;j \le m;j++)
scanf("%d",&max[i][j]);
printf("Enter
                  the
                                          matrix:\n");
                          allocation
for(i=0;i\leq n;i++)
for(j=0;j\le m;j++)
scanf("%d",&alloc[i][j]);
printf("Resource vector:");
for(i=0;i<m;i++)
scanf("%d",&total[i]);
for(i=0;i<m;i++)
avail[i]=0;
                    for(i=0;i\le n;i++)
```

```
for(j=0;j \le m;j++)
avail[j]+=alloc[i][j];
for(i=0;i<m;i++)
work[i]=avail[i];
for(j=0;j\le m;j++)
work[j]=total[j]-work[j];
for(i=0;i<n;i++)
for(j=0;j<m;j++)
need[i][j]=max[i][j]-alloc[i][j]
; A:
for(i=0;i<n;i++)
{
c=0;
for(j=0;j<m;j++)
if((need[i][j]<=work[j])&&(finish[i]=='n'))
c++;
if(c==m)
{
printf("All the resources can be allocated to Process %d", i+1);
printf("\n\nAvailable resources are:");
for(k=0;k<m;k++)
{
work[k]+=alloc[i][k];
printf("%4d",work[k]);
}
printf("\n");
finish[i]='y';
printf("\nProcess %d executed?:%c \n",i+1,finish[i]);
count++;
}
if(count!=n)
goto A;
else
printf("\n System is in safe mode");
printf("\n The given state is safe state");
getch();
```

Enter the no. of processes and resources: 43

Enter the claim matrix:

322

613

3 1 4

422

Enter the allocation matrix:

100

612

2 1 1

002

Resource vector:936

All the resources can be allocated to Process 2

Available resources are: 6 2 3

Process 2 executed?:y

All the resources can be allocated to Process 3 Available resources

are: 834

Process 3 executed?:y

All the resources can be allocated to Process 4 Available resources

are: 836

Process 4 executed?:y

All the resources can be allocated to Process 1

Available resources are: 936

Process 1 executed?:y

System is in safe mode

The given state is safe state

VIVA QUESTIONS

- 1) What is meant by deadlock?
- 2) What is safe state in banker's algorithms?
- 3) What is banker's algorithm?
- 4) What are the necessary conditions where deadlock occurs?
- 5) What are the principles and goals of protection?

EXPERIMENT.NO 10 <u>DEAD</u> <u>LOCKPREVENTION</u>

AIM: To implement deadlock prevention technique

Banker's Algorithm:

When a new process enters a system, it must declare the maximum number of instances of each resource type it needed. This number may exceed the total number of resources in the system. When the user request a set of resources, the system must determine whether the allocation of each resources will leave the system in safe state. If it will the resources are allocation; otherwise the process must wait until some other process release the resources.

DESCRIPTION:

Data structures

- n-Number of process, m-number of resource types.
- Available: Available[j]=k, k instance of resource type Rj is available.
- Max: If max[i, j]=k, Pi may request at most k instances resource Rj.

 Allocation: If Allocation [i, j]=k, Pi allocated to k instances of resource Rj Need:

 If Need[I, j]=k, Pi may need k more instances of resource type Rj,

Need[I, j]=Max[I, j]-Allocation[I, j];

Safety Algorithm

Work and Finish be the vector of length m and n respectively, Work=Available and Finish[i] =False.

Find an i such
that both Finish[i] = False
• Need <= Work

If no such I exists go to step 4.

5. work=work+Allocation, Finish[i] =True;

if Finish[1]=True for all I, then the system is in safe state

ALGORITHM:

- 1. Start the program.
- 2. Get the values of resources and processes.
- 3. Get the avail value.
- 4. After allocation find the need value.
- 5. Check whether its possible to allocate.
- 6. If it is possible then the system is in safe state.
- 7. Else system is not in safety state
- 8. Stop the process.

SOURCE CODE:

```
#include<stdio.h>
#include<conio.h
> void main()
char job[10][10];
int time[10], avail, tem[10], temp[10]; int
safe[10]; int ind=1,i,j,q,n,t;
clrscr();
printf("Enter no of jobs: ");
scanf("%d",&n);
for(i=0;i<n;i++)
printf("Enter name and time:
"); scanf("%s%d",&job[i],&time[i]);
}
printf("Enter the
                       available
resources:"); scanf("%d",&avail);
for(i=0;i<n;i++)
temp[i]=time[i];
tem[i]=i;
for(i=0;i\leq n;i++)
for(j=i+1;j < n;j++)
if(temp[i]>temp[j])
t=temp[i];
```

```
temp[i]=temp[j];
temp[j]=t; t=tem[i];
tem[i]=tem[j];
tem[j] = t;
}
 }
for(i=0;i<n;i++)
q=tem[i];
if(time[q] \le avail)
safe[ind]=tem[i];
avail=avail-tem[q];
printf("%s",job[safe[ind]])
; ind++;
}
else
printf("No safe sequence\n");
printf("Safe sequence is:");
for(i=1;i<ind; i++)
printf("%s %d\n",job[safe[i]],time[safe[i]]);
getch();
}
```

OUTPUT:

Enter no of jobs:4

Enter name and time: A 1 Enter name and time: B 4 Enter name and time: C 2 Enter name and time: D 3

Enter the available resources: 20 Safe sequence is: A 1, C 2, D 3, B 4.

EXPERIMENT.NO 11

AIM: To Write a C program to simulate disk scheduling algorithms a) FCFS b) SCAN c) C-SCAN

DESCRIPTION

One of the responsibilities of the operating system is to use the hardware efficiently. For the disk drives, meeting this responsibility entails having fast access time and large disk bandwidth. Both the access time and the bandwidth can be improved by managing the order in which disk I/O requests are serviced which is called as disk scheduling. The simplest form of disk scheduling is, of course, the first-come, first-served (FCFS) algorithm. This algorithm is intrinsically fair, but it generally does not provide the fastest service. In the SCAN algorithm, the disk arm starts at one end, and moves towards the other end, servicing requests as it reaches each cylinder, until it gets to the other end of the disk. At the other end, the direction of head movement is reversed, and servicing continues. The head continuously scans back and forth across the disk. C-SCAN is a variant of SCAN designed to provide a more uniform wait time. Like SCAN, C-SCAN moves the head from one end of the disk to the other, servicing requests along the way. When the head reaches the other end, however, it immediately returns to the beginning of the disk without servicing any requests on the return trip

PROGRAM

A) FCFS DISK SCHEDULING ALGORITHM

```
#include<stdio.h
> main()
       int t[20], n, I, j, tohm[20], tot=0; float avhm;
       clrscr();
       printf("enter
                       the
                             no.of
                                      tracks"):
       scanf("%d",&n);
       printf("enter
                            tracks
                      the
                                           be traversed");
                                     to
       for(i=2;i < n+2;i++)
               scanf("%d",&t*i
        +); for(i=1;i < n+1;i++)
               tohm[i] = t[i+1] - t[i];
               if(tohm[i] \le 0)
               tohm[i]=tohm[i]*(-1
               );
        for(i=1;i \le n+1;i++)
               tot+=tohm[i]
        ; avhm=(float)tot/n;
        printf("Tracks traversed\tDifference between tracks\n");
        for(i=1;i \le n+1;i++)
               printf("%d\t\t\d\n",t*i+,tohm*i+);
               printf("\nAverage
                                    header
                                               movements:%f",avhm);
        getch();
```



TATIOTTI
INPIII

Enter no. of tracks:9

Enter track position:55 160 184

OUTPUT

Tracks traversed

Difference between tracks

Average header movements:30.888889

B) SCAN DISK SCHEDULING ALGORITHM

```
#include<stdio.h
> main()
{
       int t[20], d[20], h, i, j, n, temp, k, atr[20], tot, p, sum=0;
        clrser();
       printf("enter the no of tracks to be traveresed");
        scanf("%d",&n);
       printf("enter the position of head");
       scanf("%d",&h);
       t[0]=0;t[1]=h;
       printf("enter the
       tracks");
       for(i=2;i\le n+2;i++)
               scanf("%d",&t[i]
        ); for(i=0;i \le n+2;i++)
     for(j=0;j \le (n+2)-i-1;j++)
     if(t[j]>t[j+1])
       1
     temp=t[j];
     t[j]=t[j+1];
     t[j+1]=temp
     }}}
     for(i=0;i<n+2;i+
     +) if(t[i]==h)
             j=i;k=i;
       p=0;
       while(t[j]!=0)
               atr[p]=t[j];
               j--, p++,
        atr[p]=t[j];
       for(p=k+1;p \le n+2;p++,k++)
               atr[p]=t[k+1]
        ; for(j=0;j\leq n+1;j++)
               if(atr[j]>atr[j+1])
                       d[j]=atr[j]-atr[j+1];
               else
                        d[j]=atr[j+1]-atr[j];
               sum+=d[i];
        }
       printf("\nAverage header movements:%f",(float)sum/n);
```

TA	TI	TT	787
1/1	130	11	

Enter no.of tracks:9
Enter track position:55

58	60	70	18
90	150	160 1	84

OUTPUT

1000 A 11 11 A	To 1 000	991 Y Y 101
Tracks traversed	Difference	between tracks

150	50
160	10
184	24
90	94
70	20
60	10
58	2
55	3
18	37

Average header movements: 27.77

C) C-SCAN DISK SCHEDULING ALGORITHM

```
#include<stdio.h
          > main()
          {
                  int\ t[20],\ d[20],\ h,\ i,\ j,\ n,\ temp,\ k,\ atr[20],\ tot,\ p,\ sum=0;
                  clrser();
                  printf("enter the no of tracks to be traveresed");
                  scanf("%d",&n);
                  printf("enter the position of head");
                  scanf("%d",&h);
                  t[0]=0;t[1]=h;
                  printf("enter
                                      total
                  tracks"); scanf("%d",&tot);
                  t[2]=tot-1;
                  printf("enter
                                              tracks");
                                     the
                  for(i=3;i\le n+2;i++)
                          scanf("%d",&t[i]
                  ); for(i=0;i\leq n+2;i++)
                          for(j=0;j \le (n+2)-i-1;j++)
                                  if(t[j]>t[j+1])
                                                                       t
                                                                       e
                                                                       m
for(i=0;i\le n+2;i++) if(t[i]==h);
       j=i;break;
```

```
1
                                                          e
1
                                                         m
=
                                                          p
t
                                                  }
         p=0;
         while(t[j]!=tot-1)
                atr[p]=t[j]
                ; j++;
                p++;
         }
         atr[p]=t[j]
         ; p++;
         i=0;
         while(p!=(n+3) && t[i]!=t[h])
         {
                atr[p] = t[i]; i++;
                p++;
         }
```

```
if(atr[j] \ge atr[j+1])
                       d[j] = atr[j] - atr[j+1];
               else
                       d[j] = atr[j+1] - atr[j];
               sum+=d[j];
       printf("total header movements%d",sum);
       printf("avg is %f",(float)sum/n);
       getch();
}
```

 $for(j=0;j\le n+2;j++)$

INPUT

Enter the track position: 55 58 60

70 18

90 150

160 184

Enter starting position: 100

OUTPUT

Tracks traversed	Difference Between tracks
150	50
160	10
184	24
18	240
55	37
58	3
60	2
70	10
90	20

Average seek time: 35.7777779